

Quantifying channels output similarity with applications to quantum control

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Abstract

In this work we aim at quantifying quantum channel output similarity. In order to achieve this, we introduce the notion of quantum channel superfidelity, which gives us an upper bound on the quantum channel fidelity. This quantity is expressed in a clear form using the Kraus representation of a quantum channel. As examples, we show potential applications of this quantity in the quantum control field.

1 Introduction

Recent applications of quantum mechanics are based on processing and transferring information encoded in quantum states. The full description of quantum information processing procedures is given in terms of quantum channels, *i.e.* completely positive, trace preserving maps on the set of quantum states.

In many areas of quantum information processing one needs to quantify the difference between ideal quantum procedure and the procedure which is performed in the laboratory. This is especially true in the situation when one deals with imperfections during the realization of experiments. These imperfections can be countered, in a quantum control setup, using various techniques, such as dynamical decoupling [1–4], sliding mode control [5] and risk sensitive quantum control [6, 7]. A different approach is to model the particular setup and optimize control pulses for a specific task in a specific

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setup [8–11]. In particular the problem of quantifying the distance between quantum channels was studied in the context of channel distinguishability.

One possible approach to quantifying the distance between two quantum channels is to consider the fidelity between Choi-Jamiołkowski states corresponding to quantum channels [?]. Another approach could involve the diamond norm [?] of quantum channels. We propose an approach which focuses on the outputs of quantum channels.

The main aim of this paper is to provide a succinct expression for the channel output similarity. As a measure of similarity we will consider the superfidelity function and define channel superfidelity. Then we will show examples of application of our results to various pairs of quantum channels. In the final part of the paper, we will study the impact of Hamiltonian errors on the channel superfidelity. First we will consider a single qubit at a finite temperature, and next we will move to an extended quantum control example.

2 Preliminaries

Henceforth we will denote the set of linear operators, transforming vectors from a finite-dimensional Hilbert space \mathcal{X} to another finite dimensional Hilbert space \mathcal{Y} by $\mathcal{L}(\mathcal{X}, \mathcal{Y})$. We put $\mathcal{L}(\mathcal{X}) = \mathcal{L}(\mathcal{X}, \mathcal{X})$. By $\mathcal{U}(\mathcal{X})$ we will denote the set of unitary operators on \mathcal{X} . Given an operator $A \in \mathcal{L}(\mathcal{X}, \mathcal{Y})$ we denote by $\|A\|_p$ its Schatten p -norm. By \bar{A} we will denote the element-wise complex conjugation of A .

2.1 Quantum states and channels

First, we introduce two basic notions: density operators and superoperators:

Definition 1 *We call an operator $\rho \in \mathcal{L}(\mathcal{X})$ a density operator iff $\rho \geq 0$ and $\text{Tr} \rho = 1$. We denote the set of all density operators on \mathcal{X} by $\mathcal{D}(\mathcal{X})$.*

From this follows that ρ is in the form $\rho = \sum_j \lambda_j |\lambda_j\rangle \langle \lambda_j|$, where λ_j and $|\lambda_j\rangle$ denote the j^{th} eigenvalue and eigenvector of ρ respectively.

Definition 2 *A superoperator is a linear mapping acting on linear operators $\mathcal{L}(\mathcal{X})$ on a finite dimensional Hilbert space \mathcal{X} and transforming them into operators on another finite dimensional Hilbert space \mathcal{Y} . Thus*

$$\Phi : \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{Y}). \quad (1)$$

Now we define the tensor product of superoperators

Definition 3 *Given superoperators*

$$\Phi_1 : \mathcal{L}(\mathcal{X}_1) \rightarrow \mathcal{L}(\mathcal{Y}_1), \Phi_2 : \mathcal{L}(\mathcal{X}_2) \rightarrow \mathcal{L}(\mathcal{Y}_2), \quad (2)$$

we define the product superoperator

$$\Phi_1 \otimes \Phi_2 : \mathcal{L}(\mathcal{X}_1 \otimes \mathcal{X}_2) \rightarrow \mathcal{L}(\mathcal{Y}_1 \otimes \mathcal{Y}_2), \quad (3)$$

to be the unique linear mapping that satisfies the equation

$$(\Phi_1 \otimes \Phi_2)(A_1 \otimes A_2) = \Phi_1(A_1) \otimes \Phi_2(A_2), \quad (4)$$

for all operators $A_1 \in \mathcal{L}(\mathcal{X}_1), A_2 \in \mathcal{L}(\mathcal{X}_2)$.

In the most general case, the evolution of a quantum system can be described using the notion of a *quantum channel* [12–14].

Definition 4 *A quantum channel is a superoperator Φ that satisfies the following restrictions:*

1. Φ is trace-preserving, i.e. $\forall A \in \mathcal{L}(\mathcal{X}) \quad \text{Tr}(\Phi(A)) = \text{Tr}(A)$,
2. Φ is completely positive, that is for every finite-dimensional Hilbert space \mathcal{Z} the product of Φ and identity mapping on $\mathcal{L}(\mathcal{Z})$ is a non-negativity preserving operation, i.e.

$$\forall \mathcal{Z} \quad \forall A \in \mathcal{L}(\mathcal{X} \otimes \mathcal{Z}) \quad A \geq 0 \Rightarrow \Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})}(A) \geq 0. \quad (5)$$

Many different representations of quantum channels can be chosen, depending on the application. In this paper we will use only the Kraus representation.

Definition 5 *The Kraus representation of a completely positive superoperator (Def. 4(2)) is given by a set of operators $K_i \in \mathcal{L}(\mathcal{X}, \mathcal{Y})$. The action of the superoperator Φ is given by:*

$$\Phi(\rho) = \sum_i K_i \rho K_i^\dagger, \quad (6)$$

This form ensures that the superoperator is completely positive. For it to be also trace-preserving we need to impose the following constraint on the Kraus operators

$$\sum_i K_i^\dagger K_i = \mathbb{1}_{\mathcal{X}}, \quad (7)$$

where $\mathbb{1}_{\mathcal{X}}$ denotes the identity operator acting on the Hilbert space \mathcal{X} .

2.2 Superfidelity

In this section we introduce the *superfidelity*, along with its properties

Definition 6 *Superfidelity of two density operators $\rho, \sigma \in \mathcal{D}(\mathcal{X})$ is given by*

$$G(\rho, \sigma) = \text{Tr}(\rho\sigma) + \sqrt{1 - \text{Tr}\rho^2}\sqrt{1 - \text{Tr}\sigma^2}. \quad (8)$$

The superfidelity is an upper bound for the fidelity function [12, 15].

Properties of the superfidelity [15] ($\rho_1, \rho_2, \rho_3, \rho_4 \in \mathcal{D}(\mathcal{X})$):

1. Bounds: $0 \leq G(\rho_1, \rho_2) \leq 1$.
2. Symmetry: $G(\rho_1, \rho_2) = G(\rho_2, \rho_1)$.
3. Unitary invariance: $G(\rho_1, \rho_2) = G(U\rho_1U^\dagger, U\rho_2U^\dagger)$, where $U \in \mathcal{U}(\mathcal{X})$.
4. Joint concavity [16]:

$$G(p\rho_1 + (1-p)\rho_2, p\rho_3 + (1-p)\rho_4) \leq pG(\rho_1, \rho_3) + (1-p)G(\rho_2, \rho_4) \quad (9)$$

for $p \in [0, 1]$.

5. Super-multiplicativity:

$$G(\rho_1 \otimes \rho_2, \rho_3 \otimes \rho_4) \geq G(\rho_1, \rho_3)G(\rho_2, \rho_4). \quad (10)$$

6. Bound for trace distance [17]

$$\frac{1}{2}\|\rho_1 - \rho_2\|_1 \geq 1 - G(\rho_1, \rho_2). \quad (11)$$

2.3 Supporting definitions

In this section we define additional operations used in our proof. We begin with the *partial trace*

Definition 7 *For all operators A, B the partial trace is a linear mapping defined as:*

$$\text{Tr}_{\mathcal{Y}} A \otimes B = A \text{Tr} B. \quad (12)$$

The extension to operators not in the tensor product form follows from linearity.

We will also need the notion of conjugate superoperator

Definition 8 *Given a quantum channel $\Phi : \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{Y})$, for every operator $A \in \mathcal{L}(\mathcal{X})$, $B \in \mathcal{L}(\mathcal{Y})$ we define the conjugate superoperator $\Phi^\dagger : \mathcal{L}(\mathcal{Y}) \rightarrow \mathcal{L}(\mathcal{X})$ as the mapping satisfying*

$$\text{Tr}(\Phi(A)B) = \text{Tr}(A\Phi^\dagger(B)). \quad (13)$$

Note, that the conjugate to completely positive superoperator is completely positive, but is not necessarily trace-preserving

Next, we will define a reshaping operation, which preserves the lexicographical order and its inverse.

Definition 9 We define the linear mapping

$$\text{res} : \mathcal{L}(\mathcal{X}, \mathcal{Y}) \rightarrow \mathcal{Y} \otimes \mathcal{X}, \quad (14)$$

for dyadic operators as

$$\text{res}(|\psi\rangle\langle\phi|) = |\psi\rangle\langle\phi|, \quad (15)$$

for $|\psi\rangle \in \mathcal{Y}$ and $|\phi\rangle \in \mathcal{X}$ and uniquely extended by linearity.

We introduce the inverse of the $\text{res}(\cdot)$

Definition 10 We define the linear mapping

$$\text{unres} : \mathcal{Y} \otimes \mathcal{X} \rightarrow \mathcal{L}(\mathcal{X}, \mathcal{Y}) \quad (16)$$

such that

$$\forall X \in \mathcal{L}(\mathcal{X}, \mathcal{Y}) \quad \text{unres}(\text{res}(X)) = X. \quad (17)$$

Remark 1 For every choice of Hilbert spaces $\mathcal{X}_1, \mathcal{X}_2, \mathcal{Y}_1$ and \mathcal{Y}_2 and every choice of operators $A \in \mathcal{L}(\mathcal{X}_1, \mathcal{Y}_1)$, $B \in \mathcal{L}(\mathcal{X}_2, \mathcal{Y}_2)$ and $X \in \mathcal{L}(\mathcal{X}_2, \mathcal{X}_1)$ it holds that:

$$(A \otimes B)\text{res}(X) = \text{res}(AXB^T) \quad (18)$$

Remark 2 For any choice of Hilbert spaces \mathcal{X} and \mathcal{Y} and any choice of $|\zeta\rangle \in \mathcal{X} \otimes \mathcal{Y}$ and $A \in \mathcal{L}(\mathcal{Y}, \mathcal{X})$ such that $|\zeta\rangle = \text{res}(A)$ it holds that

$$\text{Tr}_{\mathcal{Y}} |\zeta\rangle\langle\zeta| = \text{Tr}_{\mathcal{Y}}(\text{res}(A)\text{res}(A)^\dagger) = AA^\dagger. \quad (19)$$

Next, we introduce the *purification* of quantum states:

Definition 11 Given Hilbert spaces \mathcal{X} and \mathcal{Y} , we will call $|\zeta\rangle \in \mathcal{X} \otimes \mathcal{Y}$ a purification of $\rho \in \mathcal{D}(\mathcal{X})$ if

$$\text{Tr}_{\mathcal{Y}} |\zeta\rangle\langle\zeta| = \rho. \quad (20)$$

Theorem 1 For every choice of Hilbert spaces \mathcal{X} and \mathcal{Y} and let $|\phi\rangle, |\psi\rangle \in \mathcal{X} \otimes \mathcal{Y}$ satisfy

$$\text{Tr}_{\mathcal{Y}}(|\phi\rangle\langle\phi|) = \text{Tr}_{\mathcal{Y}}(|\psi\rangle\langle\psi|). \quad (21)$$

Then there exists a unitary operator $U \in \mathcal{U}(\mathcal{X})$ such that $|\psi\rangle = (\mathbb{1}_{\mathcal{X}} \otimes U)|\phi\rangle$

From Definition 11, Theorem 1 and Remark 2 we get that given a state $\rho \in \mathcal{D}(\mathcal{X})$ its purification $|\zeta\rangle \in \mathcal{X} \otimes \mathcal{X}$ is given by:

$$|\zeta\rangle = (\mathbb{1}_{\mathcal{X}} \otimes U)\text{res}(\sqrt{\rho}). \quad (22)$$

Verification of this equation is straightforward. First, we note that we may omit the term $(\mathbb{1}_{\mathcal{X}} \otimes U)$. Next we apply Remark 2 which allows us to show that for this choice of $|\zeta\rangle$ we get:

$$\text{Tr}_{\mathcal{Y}} |\zeta\rangle\langle\zeta| = \sqrt{\rho}(\sqrt{\rho})^\dagger = \rho. \quad (23)$$

2.4 Quantum channel fidelity

First, we introduce the *fidelity* and *channel fidelity* [?]

Definition 12 *Given two density operators ρ, σ we define the fidelity between ρ and σ as:*

$$F(\rho, \sigma) = \|\sqrt{\rho}\sqrt{\sigma}\|_1^2 = \left(\text{Tr}\sqrt{\sqrt{\rho}\sigma\sqrt{\rho}}\right)^2 \quad (24)$$

Definition 13 *Quantum channel fidelity of a channel $\Phi : \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{X})$ for some σ is defined as:*

$$F_{\text{ch}}(\Phi; \sigma) = \inf F(\xi, (\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(\xi)), \quad (25)$$

where the infimum is over all Hilbert spaces \mathcal{Z} and all $\xi \in \mathcal{D}(\mathcal{X} \otimes \mathcal{Z})$ such that $\text{Tr}_{\mathcal{Z}}\xi = \sigma$

It can be shown [18] that this infimum is independent of ξ and is given by

$$F_{\text{ch}}(\Phi; \sigma) = \sum_i |\text{Tr}(\sigma K_i)|^2, \quad (26)$$

where K_i form the Kraus representation of Φ .

3 Our results

In this section we present our main theorem and its proof. In the second subsection we present a quantum circuit that allows one to measure the quantum channel superfidelity without performing full state tomography.

3.1 Theorem and proof

Definition 14 *Consider two quantum channels $\Phi, \Psi : \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{X})$ and a density operator $\sigma \in \mathcal{D}(\mathcal{X})$. We define the quantum channel superfidelity to be:*

$$G_{\text{ch}}(\Phi, \Psi; \sigma) = \inf G((\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(\xi), (\Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(\xi)), \quad (27)$$

where the infimum is over all Hilbert spaces \mathcal{Z} and over all purifications $\xi = |\zeta\rangle\langle\zeta| \in \mathcal{D}(\mathcal{X} \otimes \mathcal{Z})$ of σ .

The channel superfidelity $G_{\text{ch}}(\Phi, \Psi; \sigma)$ places a lower bound on the output superfidelity of two quantum channels in the case of the same input states. Henceforth, where unambiguous, we will write the channel superfidelity as G_{ch} .

Theorem 2 *Given quantum channels $\Phi, \Psi : \mathcal{L}(\mathcal{X}) \rightarrow \mathcal{L}(\mathcal{X})$ with Kraus forms given by the sets $\{K_i\}_i$ and $\{L_j\}_j$ respectively the quantum channel superfidelity is given by:*

$$G_{\text{ch}} = \sum_{i,j} |\text{Tr} \sigma K_i^\dagger L_j|^2 + \sqrt{1 - \sum_{i,j} |\text{Tr} \sigma K_i^\dagger K_j|^2} \times \sqrt{1 - \sum_{i,j} |\text{Tr} \sigma L_i^\dagger L_j|^2}. \quad (28)$$

Proof. As we limit ourselves only to pure states ξ , in order to calculate the superfidelity, we need to compute the following quantities:

$$\begin{aligned} & \text{Tr} (\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})}) (|\zeta\rangle \langle \zeta|) (\Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})}) (|\zeta\rangle \langle \zeta|), \\ & \text{Tr} [(\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})}) (|\zeta\rangle \langle \zeta|)]^2, \\ & \text{Tr} [(\Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})}) (|\zeta\rangle \langle \zeta|)]^2. \end{aligned} \quad (29)$$

As the general idea is shared between all of these quantities, we will show here the calculation for the first one. We get

$$\begin{aligned} & \text{Tr} (\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})}) (|\zeta\rangle \langle \zeta|) (\Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})}) (|\zeta\rangle \langle \zeta|) = \\ & \text{Tr} |\zeta\rangle \langle \zeta| \left(\Phi^\dagger \circ \Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})} \right) (|\zeta\rangle \langle \zeta|) = \\ & \langle \zeta| \left(\Phi^\dagger \circ \Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})} \right) (|\zeta\rangle \langle \zeta|) |\zeta\rangle, \end{aligned} \quad (30)$$

where the first equality follows from the definition of the conjugate superoperator. The Kraus form of the superoperator $\Phi^\dagger \circ \Psi$ is given by the set $\{K_i^\dagger L_j\}_{i,j}$. Now, we may write $|\zeta\rangle$ as

$$|\zeta\rangle = \text{res}(\sqrt{\sigma}U), \quad (31)$$

for some $U \in \mathcal{L}(\mathcal{Z}, \mathcal{X})$ such that UU^\dagger is a projector on the image of σ . We obtain:

$$\begin{aligned} & \langle \zeta| \left(\Phi^\dagger \circ \Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})} \right) (|\zeta\rangle \langle \zeta|) |\zeta\rangle = \sum_{i,j} |\langle \zeta| (K_i^\dagger L_j) \otimes \mathbb{1}_{\mathcal{Z}} |\zeta\rangle|^2 = \\ & \sum_{i,j} |\text{res}(\sqrt{\sigma}U)^\dagger \left((K_i^\dagger L_j) \otimes \mathbb{1}_{\mathcal{Z}} \right) \text{res}(\sqrt{\sigma}U)|^2 = \\ & \sum_{i,j} |\text{Tr} U^\dagger \sqrt{\sigma} K_i^\dagger L_j \sqrt{\sigma} U|^2 = \sum_{i,j} |\text{Tr} \sigma K_i^\dagger L_j|^2. \end{aligned} \quad (32)$$

This quantity is independent of the particular purification of σ . Following the same path for the other two quantities shown in Eq (29), we recover the expression for the channel superfidelity from Eq (28). \blacksquare

Since the superfidelity is an upper bound for the fidelity function, we obtain, the following inequality:

$$G_{\text{ch}}(\Phi, \Psi; \sigma) \geq F(\Phi, \Psi; \sigma), \quad (33)$$

where $F(\Phi, \Psi; \sigma) = \inf_{\mathcal{Z}, \xi} F((\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(\xi), (\Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(\xi))$ and the infimum is over all Hilbert spaces \mathcal{Z} and all $\xi \in \mathcal{D}(\mathcal{X} \otimes \mathcal{Z})$ such that $\text{Tr}_{\mathcal{Z}} \xi = \sigma$.

The following simple corollaries are easily derived from Theorem 2.

Corollary 1 *Given a quantum channel Φ , the superfidelity between its input and output reduces to the channel fidelity of Φ :*

$$G_{\text{ch}}(\Phi, \mathbb{1}; \sigma) = F_{\text{ch}}(\Phi; \sigma) \quad (34)$$

Proof. Assume Φ has the Kraus form $\{K_i\}_i$. Substituting the identity for L_j in Eq. (27) we recover Eq. (25) which completes the proof. ■

Corollary 2 *If Φ is a unitary channel i. e. $\Phi(\rho) = U\rho U^\dagger$ for any $U \in \mathcal{U}(\mathcal{X})$ and $\Psi' : \rho \mapsto U^\dagger \Psi(\rho) U$, where Ψ is an arbitrary quantum channel then*

$$G_{\text{ch}}(\Phi, \Psi; \sigma) = F_{\text{ch}}(\Psi'; \sigma). \quad (35)$$

Proof. If Φ is a unitary channel, then the second term in Eq. (28) vanishes. Let us assume that Ψ has a Kraus form $\{L_j\}_j$. We get $G_{\text{ch}} = \sum_j |\text{Tr} \sigma U^\dagger L_j|^2$.

The Kraus form of the channel Ψ' is given by the set $\{U^\dagger L_j : K_j \in \mathcal{L}(\mathcal{X})\}$. Using this in Eq. (25) we get $F_{\text{ch}}(\Psi'; \sigma) = \sum_j |\text{Tr} \sigma U^\dagger L_j|^2$. This completes the proof. ■

Corollary 3 *If $\sigma \in \mathcal{D}(\mathcal{X})$ is a pure state i.e. $\sigma = |\psi\rangle\langle\psi|$ then*

$$G_{\text{ch}}(\Phi, \Psi; \sigma) = G(\Phi(\sigma), \Psi(\sigma)). \quad (36)$$

Proof. Let us only focus on the first terms in Eq. (8) and Eq. (28). We will denote these terms T and T_{ch} respectively. Let us assume that channels Φ and Ψ have Kraus forms $\{K_i\}_i$ and $\{L_j\}_j$ respectively. We get:

$$\begin{aligned} T &= \text{Tr} \sum_{ij} K_i |\psi\rangle\langle\psi| K_i^\dagger L_j |\psi\rangle\langle\psi| L_j^\dagger \\ &= \sum_{ij} \langle\psi| K_i^\dagger L_j |\psi\rangle \langle\psi| L_j^\dagger K_i |\psi\rangle = \sum_{ij} |\langle\psi| K_i^\dagger L_j |\psi\rangle|^2 = T_{\text{ch}}. \end{aligned} \quad (37)$$

Performing similar calculations for other terms, we recover Eq. (36) ■

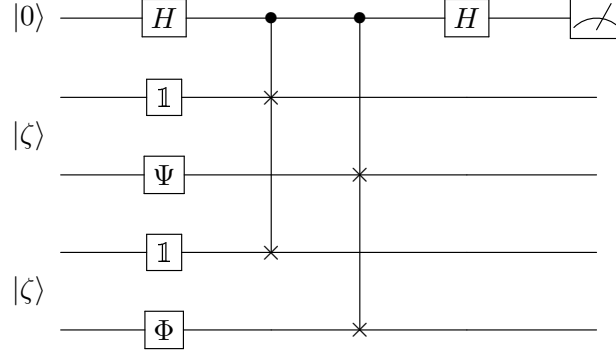


Figure 1: Quantum circuit for measuring $\text{Tr}(\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(|\zeta\rangle\langle\zeta|)(\Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(|\zeta\rangle\langle\zeta|) = 2p_0 - 1$, where p_0 is the probability of finding the top qubit in state $|0\rangle$ [19]. This allows direct estimation of the channel superfidelity.

3.2 Quantum circuit for measuring channel superfidelity

Using the quantum circuit shown in Fig. 1 we can measure the quantum channel superfidelity in an experimental setup. This setup allows us to estimate the $\text{Tr}(\Phi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(|\zeta\rangle\langle\zeta|)(\Psi \otimes \mathbb{1}_{\mathcal{L}(\mathcal{Z})})(|\zeta\rangle\langle\zeta|) = 2p_0 - 1$, where p_0 is the probability of finding the top qubit in the state $|0\rangle$. Modifying the circuit appropriately, we can measure all the quantities shown in Eq. (29).

Note that this approach is far simpler, compared to estimating the channel fidelity which would require us to perform full state tomography. Furthermore, analytical calculations involving fidelity get cumbersome quickly, as it requires calculating expressions of the form $\|\sqrt{\Phi(\sigma)}\sqrt{\Psi(\sigma)}\|_1$.

4 Simple examples

In this section we provide a number of examples of the application of Theorem 2.

4.1 Erasure channel

Definition 15 *Given a quantum state $\xi \in \mathcal{D}(\mathcal{X})$ the erasure channel is given by:*

$$\Phi(A) = \xi, \quad (38)$$

for any A in $\mathcal{L}(\mathcal{X})$. The Kraus form of this channel is given by the set $\{K_{ij} : K_{ij} = \sqrt{\lambda_i}|\lambda_i\rangle\langle j|\}_{ij}$, λ_i and $|\lambda_i\rangle$ denote the i^{th} eigenvalue and the corresponding eigenvector of ξ .

Let us consider the superfidelity between the erasure channel Φ and a unitary channel $\Psi : \sigma \mapsto U\sigma U^\dagger$ for some $U \in \mathcal{U}(\mathcal{X})$. We note that the

second term in Eq. (28) vanishes. What remains is:

$$\begin{aligned} G_{\text{ch}} &= \sum_{ij} \lambda_i |\langle j | \sigma U^\dagger | \lambda_i \rangle|^2 = \sum_j \langle j | \sigma U^\dagger \left(\sum_i \lambda_i | \lambda_i \rangle \langle \lambda_i | \right) U \sigma | j \rangle \\ &= \text{Tr} \sigma^2 \Psi^\dagger(\xi) \leq \sum_i \lambda_i^\downarrow \mu_i^\downarrow, \end{aligned} \quad (39)$$

where μ_i^\downarrow and λ_i^\downarrow denote the eigenvalues of σ and ξ respectively, sorted in a descending order. The last inequality follows from von Neumann's trace inequality [20].

4.2 Sensitivity to channel error

Consider a quantum channel Φ with the Kraus form $\{K_i\}$ and a quantum channel $\Psi : \rho \mapsto U_\epsilon \Phi(\rho) U_\epsilon^\dagger$, where $U_\epsilon = \exp(-i\epsilon H) \in \mathcal{U}(\mathcal{X})$. We get:

$$G_{\text{ch}}(\Phi, \Psi; \sigma) = 1 + \sum_{ij} |\text{Tr} \sigma K_i^\dagger U_\epsilon K_j|^2 - \sum_{ij} |\text{Tr} \sigma K_i^\dagger K_j|^2. \quad (40)$$

Now, we concentrate on the change of the quantum channel superfidelity under the change of ϵ . As we are interested only in small values of ϵ , we expand Eq. (40) up to the linear term in the Taylor series. For small values of ϵ we get:

$$G_{\text{ch}} \approx 1 - 2\epsilon \sum_{ij} \Im \text{Tr} \sigma K_i^\dagger H K_j \overline{\text{Tr} \sigma K_i^\dagger K_j}. \quad (41)$$

Note that this depends on the value of the observable H of the operator $K_j \sigma K_i^\dagger$.

5 Sensitivity to Hamiltonian parameters

In this section we will show how the channel superfidelity is affected by errors in the system Hamiltonian parameters. First, we will show analytical results for a single qubit system at a finite temperature. Next, we show numerical results for a simple, three-qubit spin chain.

5.1 Single qubit at a finite temperature

A single qubit at a finite temperature is described by the master equation

$$\begin{aligned} \dot{\rho}(t) &= -i \frac{\Omega + \epsilon}{2} [\sigma_z, \rho(t)] + \gamma_+ \left(\sigma_- \rho(t) \sigma_+ - \frac{1}{2} \{ \sigma_+ \sigma_-, \rho(t) \} \right) + \\ &\quad + \gamma_- \left(\sigma_+ \rho(t) \sigma_- - \frac{1}{2} \{ \sigma_- \sigma_+, \rho(t) \} \right), \end{aligned} \quad (42)$$

where $\sigma_+ = |1\rangle\langle 0|$, $\sigma_- = \sigma_+^\dagger$ and ϵ is the error in Ω . Our goal is to calculate the quantum channel superfidelity between the case when there is no error in Ω , i.e. $\epsilon = 0$ and the case with error in Ω . Henceforth, we will assume $\gamma_- = \gamma_+ = 1$ for clarity.

For a given time T , Eq. (42) may be rewritten as

$$\rho(T) = \Phi_T^\epsilon(\rho(0)), \quad (43)$$

where Φ_T^ϵ is a quantum channel in the quantum dynamical semigroup. A natural representation $M_{\Phi_T^\epsilon}$ for the channel Φ_T^ϵ may be found as [21]:

$$M_{\Phi_T^\epsilon} = e^{-FT}, \quad (44)$$

where

$$\begin{aligned} F = & -i \frac{\Omega + \epsilon}{2} (\mathbb{1} \otimes \sigma_z - \sigma_z \otimes \mathbb{1}) - \sigma_- \otimes \sigma_- - \sigma_+ \otimes \sigma_+ + \\ & + \frac{1}{2} (\sigma_+ \sigma_- \otimes \mathbb{1} + \sigma_- \sigma_+ \otimes \mathbb{1} + \mathbb{1} \otimes \sigma_+ \sigma_- + \mathbb{1} \otimes \sigma_- \sigma_+). \end{aligned} \quad (45)$$

In this representation we may rewrite Eq. (43) as

$$\text{res}(\rho(T)) = M_{\Phi_T^\epsilon} \text{res}(\rho(0)). \quad (46)$$

The Choi-Jamiołkowski representation of the channel Φ_T is given by $D_{\Phi_T^\epsilon} = (M_{\Phi_T^\epsilon})^R$. Here, M^R denotes the *reshuffle* operation on matrix M [12]. Now, it is simple to find the Kraus form of the channel Φ_T^ϵ . The Kraus operators are related to the eigenvalues λ_i and eigenvectors $|\lambda_i\rangle$ of $D_{\Phi_T^\epsilon}$ in the following manner:

$$K_i^{\Phi_T^\epsilon} = \sqrt{\lambda_i} \text{unres}(|\lambda_i\rangle). \quad (47)$$

Inserting these Kraus operators into Eq. (28) we get

$$G_{\text{ch}}(\Phi_T^0, \Phi_T^\epsilon; \rho) = 1 - 2e^{-2T}(1 - \cos \epsilon T) \rho_{00}(0) \rho_{11}(0), \quad (48)$$

where $\rho_{ii}(0) = \langle i | \rho(0) | i \rangle$. Note that, we get $G_{\text{ch}} = 1$ in two cases. First, for large T and second when $\epsilon T = \frac{\pi}{2}$. As we are mainly interested in small values of ϵ , we expand $\cos \epsilon T$ up to the second term in the Taylor series. We get:

$$G_{\text{ch}}(\Phi_T^0, \Phi_T^\epsilon; \rho) \approx 1 - \epsilon^2 T^2 e^{-2T} \rho_{00}(0) \rho_{11}(0). \quad (49)$$

In this setup the channel superfidelity has a quadratic dependence on the error parameter ϵ . This should be compared with the results in Sec. 4.2.

5.2 Quantum control example

In this section, we consider a three-qubit spin chain with dephasing interactions with the environment. We will consider piecewise constant control pulses. The time evolution of the system is governed by the equation:

$$\dot{\rho}(t) = -i[H, \rho(t)] + \gamma(\sigma_z \rho(t) \sigma_z - \rho(t)), \quad (50)$$

where $H = H_d + H_c$. Here H_d is the drift term of the Hamiltonian given by

$$H_d = J \sum_{i=1}^2 \sum_{\alpha \in \{x,y,z\}} \sigma_{\alpha}^i \sigma_{\alpha}^{i+1}, \quad (51)$$

where σ_{α}^i denotes σ_{α} acting on site i . We set the control Hamiltonian H_c to:

$$H_c = \sum_i^N h_x(t_i) \sigma_x^1 + h_y(t_i) \sigma_y^1, \quad (52)$$

where $h_x(t_i)$ and $h_y(t_i)$ denote the control pulses in the time interval t_i . We set the target to be

$$U_T = \mathbb{1} \otimes \mathbb{1} \otimes \sigma_x, \quad (53)$$

i.e. a NOT gate on the third qubit. We fixed the number of time intervals $N = 64$, the total evolution time $T = 6.1$ and the maximum amplitude of a single control pulse $\forall k \in x, y \max(|h_k|) = 10$.

First, we optimize control pulses for the system, such that we achieve a high fidelity of the gate U_T . Next, to each control pulse we add a noise term h_{ϵ} which has a normal distribution, $h_{\epsilon} = N(0, s)$. Fig. 2 shows the change of G_{ch} as a function of the standard deviation s . We have conducted 100 simulations for each value of s . As expected, the quantum channel superfidelity decreases slowly for low values of s . After a certain value the decrease becomes rapid. As values of s increase, the minimum and maximum achieved fidelity diverge rapidly. This is represented by the shaded area in Fig. 2. We can approximate the average value of the channel fidelity as $\langle G_{\text{ch}} \rangle \approx 1 - cs^2$. Fitting this function to the curve shown in Fig. 2b gives a relative error which is less than 0.5%.

6 Conclusions

We have studied the superfidelity of a quantum channel. This quantity allows us to provide an upper bound on the fidelity of the output of two quantum channels. We shown an example of application of this quantity to a unitary and an erasure channel. The obtained superfidelity can be easily limited from above by the product w eigenvalues of the input state σ and the result of the erasure channel ξ .

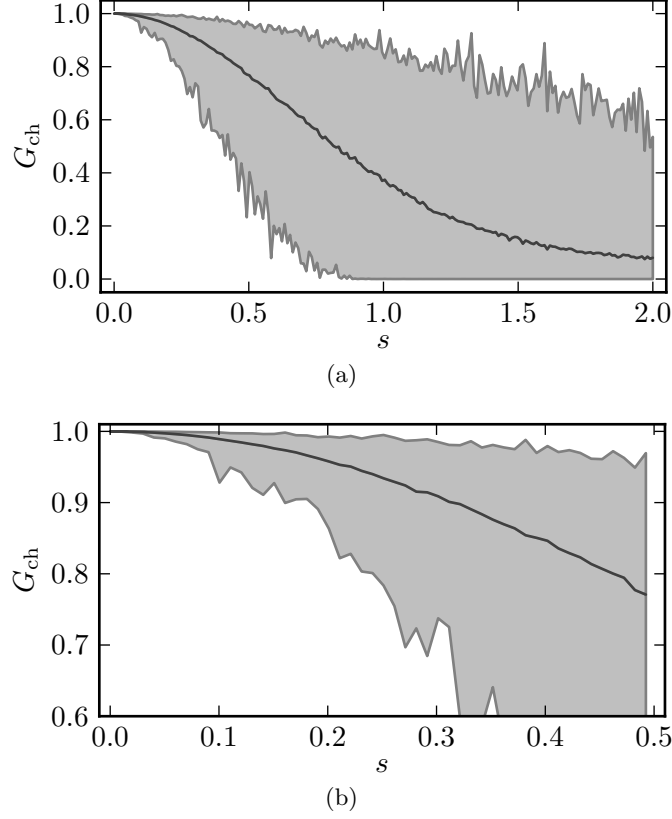


Figure 2: Quantum channel superfidelity as a function of noise in the system's control pulses. The shaded area represents the range of the achieved channel superfidelity, the black line is the average channel superfidelity.

Furthermore, as shown in our examples, the quantum channel superfidelity may have potential applications in quantum control theory as an easy to compute figure of merit of quantum operations. In a simple setup, where the desired quantum channel is changed by a unitary transformation $U_\epsilon = \exp(-i\epsilon H)$ we get a linear decrease of channel superfidelity on the noise parameter ϵ . On the other hand, when we introduce the noise as a control error in a single qubit quantum control setup, we get a quadratic dependence on the noise parameter.

Finally, we shown numerical results for a more complicated system. We calculated the quantum channel superfidelity for a three-qubit quantum control setup. First we found control pulses which achieve a high fidelity of the desired quantum operation, next we introduced Gaussian noise in the control pulses. Our results show, that the quantum channel superfidelity stayed high for a wide range of the noise strength.

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